# Cryptography and Architectures for Computer Security 

Exam Code: 095947 (old 090959), A.Y. 2018-2019, Semester: 2
Prof. G. Pelosi
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Name: . Surname: $\qquad$
Student ID: Signature: $\qquad$
Time: 2h, 15min. Use of textbooks, notes, or Internet connected devices (including smartphones) is not allowed. The usage of simple calculators is allowed. Prior to turn in your paper, write your name on any additional sheet and sign it.

## Question 1 [5 pts]

Suppose you are in charge of security for a major web site, and you are considering what would happen if an attacker stole your database of usernames and passwords. You have already implemented a basic defense: instead of storing the plaintext passwords, you store their SHA-256 hashes. Your threat model assumes that the attacker can carry out 4 million SHA-256 hashes per second. His goal is to recover as many plaintext passwords as possible from the information in the stolen database. Valid passwords for your site may contain only characters $\mathrm{a}-\mathrm{z}, \mathrm{A}-\mathrm{Z}$, and $0-9$, and are exactly 8 characters long.
(a) Given the hash of a single password, how many hours would it take for the attacker to crack a single password by brute force, on average?
(b) Describe an attack based on the use of rainbow tables, specifying what they are and the computation-memory tradeoffs that improve on the brute force approach.
(c) Describe possible improvements to the basic approach of storing usernames and hashes (where hash $=$ SHA-256(password) for the users of the website at hand, with the aim to provide strong protection agaist attacks based on the use of rainbow tables.

## Question 2 [3 pts]

A programmer wants to use CBC in order to protect both the integrity and confidentiality of network packets. She attaches a block of zero bits $M_{n+1}$ to the end of the plaintext $M_{1}\left\|M_{2}\right\| \ldots \| M_{n}$ as redundancy, then encrypts with CBC. At the receiving end, she verifies that the added redundant bits are still all zero after CBC decryption. Does this test ensure the integrity of the transferred message? Justify your answer.

## Question 3 [8 pts]

Consider the finite field $\mathbb{F}_{3^{4}}$.
(a) Establish if $f(x)=x^{4}-x-1 \in \mathbb{F}_{3}[x]$, and $g(x)=x^{4}-x^{2}-1 \in \mathbb{F}_{3}[x]$ are irreducible polynomials or not. Establish also if $f(x)$ and $g(x)$ are primitive polynomials. Justify your answer.
(b) Exhibit the number fo generators in $\mathbb{F}_{3^{4}}$, and the number of both monic irreducible polynomials and monic primitive polynomials having degree 4 and coefficients in $\mathbb{F}_{3}[x]$.
(c) Explain what is the hard problem in the multiplicative group of a finite field suitable to set up a cryptosystem and describe at least one algorithm (except brute-forcing) to solve it, specifying its computational complexity.

## Question 4 [8 pts]

Consider an El-Gamal signature scheme based on the cyclic group $G=\left(\mathbb{F}_{26}^{*}, \cdot\right)$, with order $n=|G|=63$. Let $P(x)=x^{6}+x^{5}+x^{4}+x+1 \in \mathbb{F}_{2}[x]$ be the generating polynomial of the field $\mathbb{F}_{2^{6}} \cong \mathbb{F}_{2}(\alpha), P(\alpha)=0$, $\alpha \in \mathbb{F}_{2^{6}} \backslash \mathbb{F}_{2}$, and assume to employ a hash function ( $h:\{0,1\}^{*} \mapsto \mathbb{Z}_{n}$ ) that maps its input binary sequences in their corresponding decimal values modulo $n$.
Given the key-pair:

$$
k_{\text {priv }}=\left\langle s \in \mathbb{Z}_{63}\right\rangle=\langle 11\rangle, \quad k_{\text {pub }}=\left\langle n, \alpha, \alpha^{s}\right\rangle=\left\langle 63, \alpha, \alpha^{11} \equiv \alpha^{3}+1=(001001)\right\rangle
$$

(a) Check if the following signatures are correct:
(i) $S_{1}=\left\langle m_{1}, \gamma, \delta\right\rangle=\langle(100000),(010000), 27\rangle$
(ii) $S_{2}=\left\langle m_{2}, \gamma, \delta\right\rangle=\langle(111111),(001000), 0\rangle$
(b) In a practical implementation of the ElGamal signature scheme based on the arithmetic of a finite field, what are the criteria to select the cryptosystem parameters and establish its (mathematical) security level?

Question 5 [12 pts]
(a) Apply the Pollard's $\rho$ method to factorize the RSA modulus $n=p \cdot q=713$.

Assume $f(x)=x^{2}+1 \bmod n$ as "random-walking" function. Show every step of the computation. (As a backup alternative, apply a "trivial division" strategy).
(b) Describe the Miller-Rabin primality test, and show its application to the greatest factor of $n$, with base $a=3$.
(c) Choose an admissible public exponent $e$ between the values $e=11_{\text {dec }}$ and $e=13_{\text {dec }}$ and compute the value of the corresponding RSA private key $k_{\text {priv }}=(p, q, \varphi(n), d)$. Show every step of the computation.
(d) Sign the message $m=100_{\text {dec }} \in \mathbb{Z}_{n}$ (without employing any padding scheme) through applying the CRT. Describe each step of the procedure.
(e) Assume to work into the Montgomery domain: $\left(\widetilde{\mathbb{Z}}_{p},+, \times\right), p=13$

- Show the definition of the Montgomery Multiplication and the smallest admissible value for the Montgomery Radix: $R$
- Show the high-level pseudo-code to implement the Montgomery Reduction procedure MRed (...), and prove the correctness of the algorithm.
- Compute a pair of integer values $R^{\prime}, p^{\prime}$ that satisfy the relation: $\operatorname{gcd}(R, p)=R R^{\prime}-p p^{\prime}=1$.
- Compute the Montgomery multiplication $C=A \times B \bmod p$, where $A=11_{\text {dec }}$ and $B=4_{\text {dec }}$ are values in the Montgomery domain, assuming a binary encoding of the operands

